

1.7 Recursion

Objectives

- To learn the concept and usage of Recursion in C
- Examples of Recursion in C

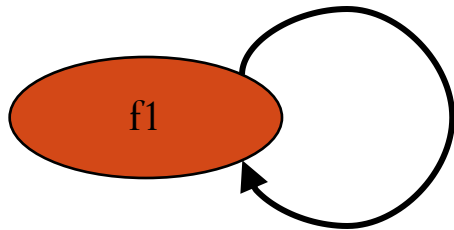
What is recursion?

- Sometimes, the best way to solve a problem is by solving a smaller version of the exact same problem first
- Recursion is a technique that solves a problem by solving a smaller problem of the same type
- a function that calls itself

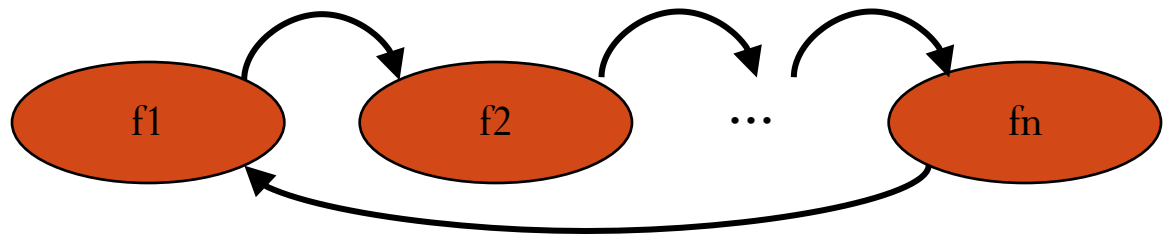
Directly or

Indirectly (a function that is part of a cycle in the sequence of function calls.)

Pictorial representation of direct and indirect recursive calls



Direct recursive call



Indirect recursive call

Syntax

```
function_name(parameter list)
{
    ...
    // 'c' statements
    ...
    function_name(parameter values) // recursive call
    ...
}
```

Problems defined recursively

- There are many problems whose solution can be defined recursively

Example: n factorial

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ (n-1)! * n & \text{if } n > 0 \end{cases} \quad (\text{recursive solution})$$

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ 1 * 2 * 3 * \dots * (n-1) * n & \text{if } n > 0 \end{cases} \quad \begin{array}{l} (\text{closed form solution}) \\ (\text{also called as iterative method}) \end{array}$$

Coding the factorial function

- Iterative implementation

```
int Factorial(int n)
{
    int fact = 1;

    for(int count = 2; count <= n; count++)
        fact = fact * count;

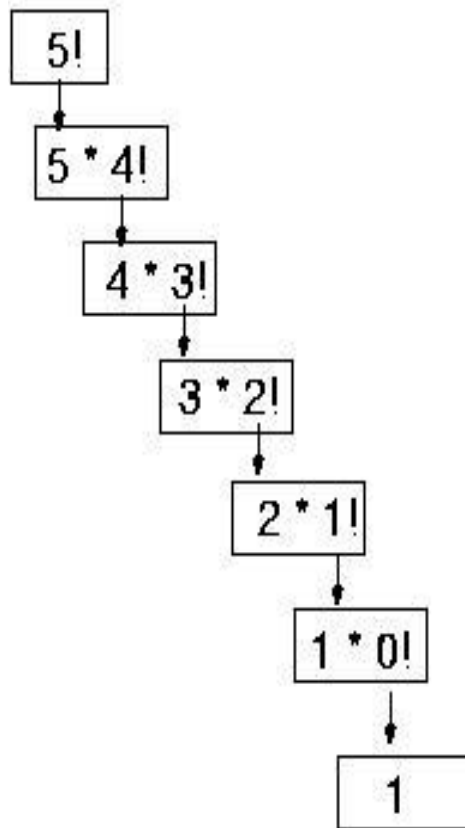
    return fact;
}
```

Coding the factorial function (An Example of Recursive Call)

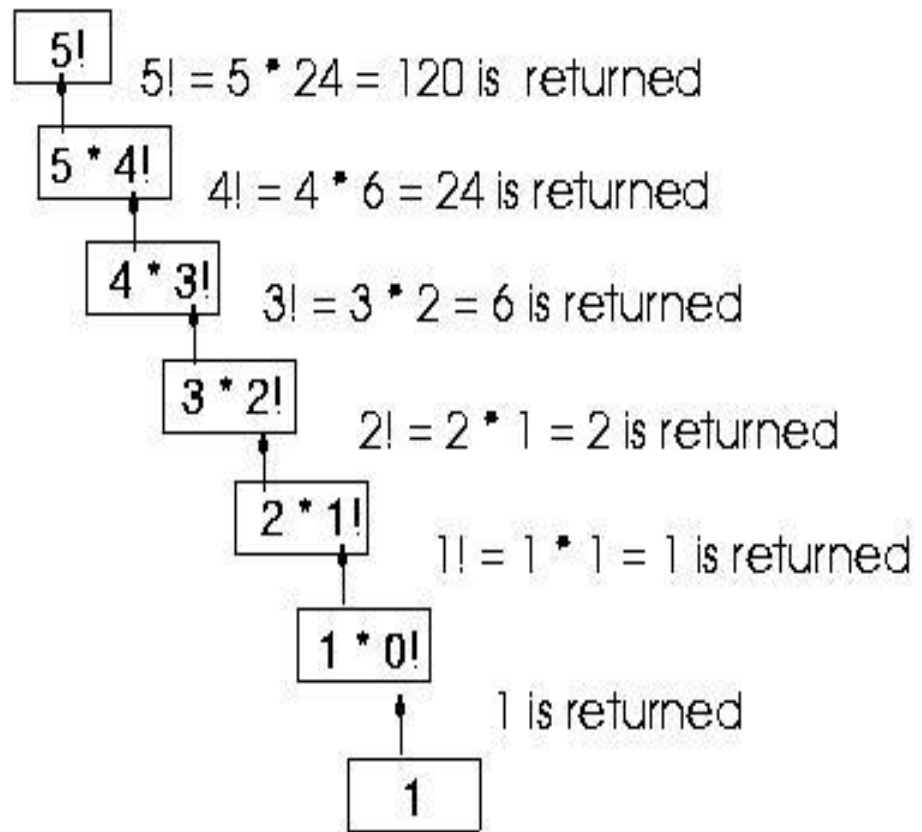
- Recursive implementation

```
int Factorial(int n)
{
    if (n==0) // base case
        return 1;
    else
        return n * Factorial(n-1);
}
```


Coding the factorial function (cont.)



Final value = 120



$x = 3$
 $y = ?$ $2 * f(2)$
call $f(2)$

push copy of f

$x = 2$
 $y = ?$ $2 * f(1)$
call $f(1)$

push copy of f

$x = 1$
 $y = ?$ $2 * f(1)$
call $f(0)$

push copy of f

$x = 0$
 $y = ?$
return $\textcircled{1}$ $= f(0)$

pop copy of f

$y = 2 * 1 = 2$
return $y + 1 = \textcircled{3}$ $= f(1)$

pop copy of f

$y = 2 * 3 = 6$
return $y + 1 = \textcircled{7}$ $= f(2)$

pop copy of f

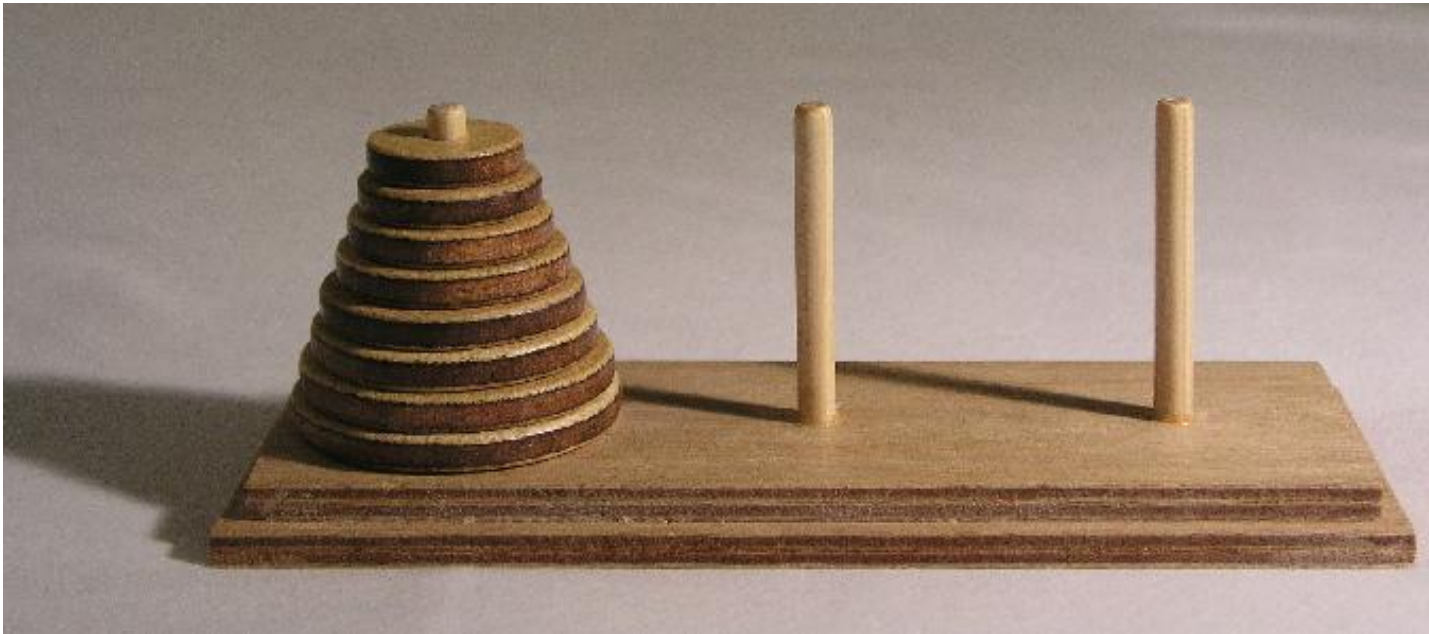
$y = 2 * 7 = 14$
return $y + 1 = \textcircled{15}$ $= f(3)$

pop copy of f

value returned by call is 15

More Interesting Example

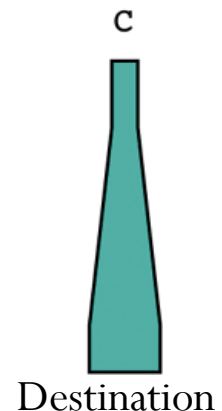
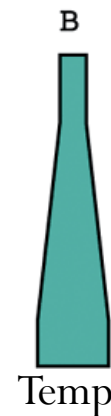
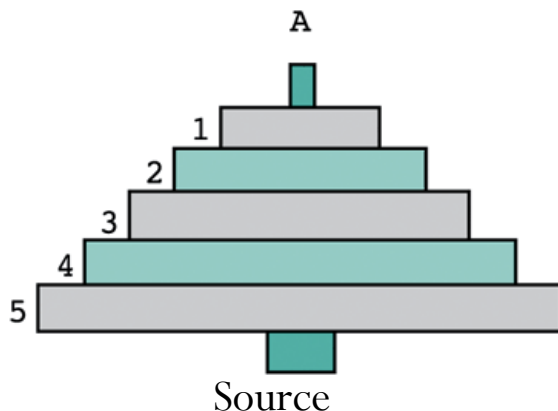
Towers of Hanoi



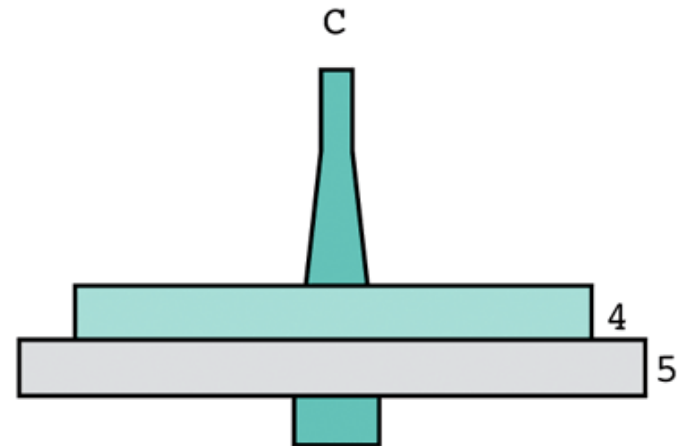
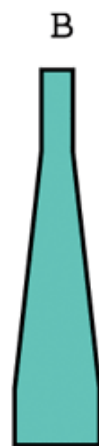
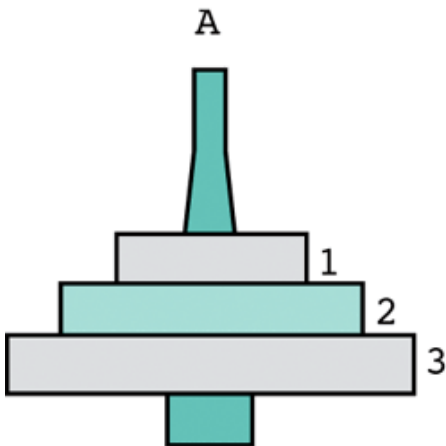
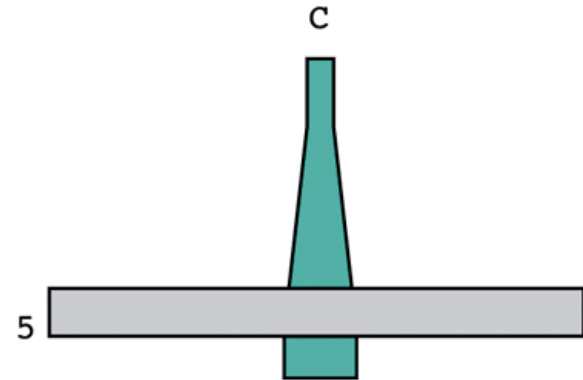
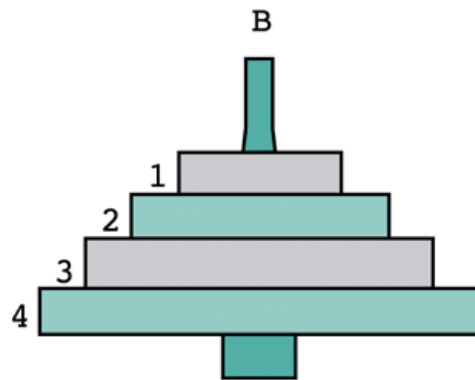
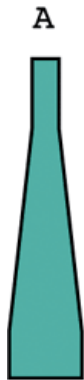
- Move stack of disks from one peg to another
- Move one disk at a time
- Larger disk may never be on top of smaller disk

A Classical Case: Towers of Hanoi

- The towers of Hanoi problem involves moving a number of disks (in different sizes) from one tower (or called “peg”) to another.
 - The constraint is that the larger disk can never be placed on top of a smaller disk.
 - Only one disk can be moved at each time
 - Assume there are three towers available.



A Classical Case: Towers of Hanoi



A Classical Case: Towers of Hanoi

- This problem can be solved easily by recursion.
- Algorithm:
 - if n is 1 then
 - move disk 1 from the source tower to the destination tower
 - else
 - 1. move $n-1$ disks from the source tower to the temp tower.
 - 2. move disk n from the source tower to the destination tower.
 - 3. move $n-1$ disks from the temp tower to the source tower.

Tower of Hanoi Program

```
#include <stdio.h>
```

```
void move (int n, int a, int  
           c, int b);
```

```
int main() {  
    int disks;  
    printf ("How many disks?");  
    scanf ("%d", &disks);  
  
    move (disks, 1, 3, 2);  
  
    return 0;  
} // main
```

```
/* PRE:  n >= 0. Disks are arranged  
         small to large on the pegs a, b,  
         and c. At least n disks on peg  
         a. No disk on b or c is smaller  
         than the top n disks of a.
```

```
POST: The n disks have been moved  
       from a to c. Small to large  
       order is preserved. Other disks  
       on a, b, c are undisturbed. */
```

```
void move (int n, int a, int c, int  
           b) {  
    if (n > 0)  
    {  
        move (n-1, a, b, c);  
        printf ("Move one disk  
                from %d to %d\n", a, c);  
        move (n-1, c, a, b);  
    }  
}
```

- Is pre-condition satisfied before this call to **move**?

Tower of Hanoi Program

```
#include <stdio.h>

void move (int n, int a, int c, int b);

int main() {
    int disks;
    printf ("How many disks?");
    scanf ("%d", &disks);

    move (disks, 1, 3, 2);
```

- If pre-condition is satisfied here, is it still satisfied here?

/* PRE: $n \geq 0$. Disks are arranged small to large on the pegs a, b, and c. At least n disks on peg a. No disk on b or c is smaller than the top n disks of a.

POST: The n disks have been moved from a to c. Small to large order is preserved. Other disks on a, b, c are undisturbed. */

```
void move (int n, int a, int c, int b) {
    if (n > 0)
    {
        move (n-1, a, b, c);
        printf ("Move one disk from %d to %d\n", a, c);
        move (n-1, b, c, a);
    } // if (n > 0)

    return;
} // move
```

And here?

Tower of Hanoi Program

```
#include <stdio.h>

void move (int n, int a, int c, int b);

int main() {
    int disks;
    printf ("How many disks?");
    scanf ("%d", &disks);

    move (disks, 1, 3, 2);

    return 0;
} // main
```

If pre-condition is true and if $n = 1$, does **move** satisfy the post-condition?

/* PRE: $n \geq 0$. Disks are arranged small to large on the pegs a, b, and c. At least n disks on peg a. No disk on b or c is smaller than the top n disks of a.

POST: The n disks have been moved from a to c. Small to large order is preserved. Other disks on a, b, c are undisturbed. */

```
void move (int n, int a, int c, int b) {
    if (n > 0)
    {
        move (n-1, a, b, c);
        printf ("Move one disk from %d to %d\n", a, c);
        move (n-1, b, c, a);
    } // if (n > 0)

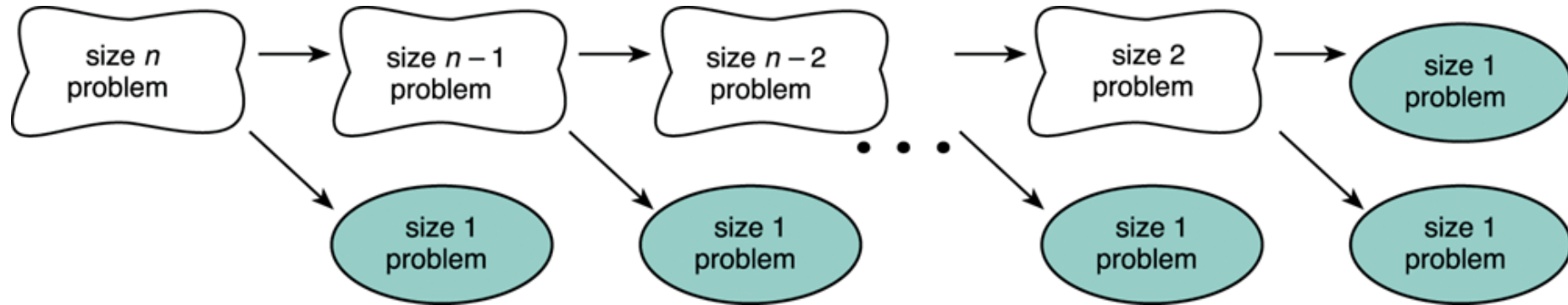
    return;
} // move
```

Can we reason that this program correctly plays Tower of Hanoi?

Problems Suitable for Recursive Functions

- One or more simple cases of the problem have a straightforward solution.
- The other cases can be redefined in terms of problems that are closer to the simple cases.
- The problem can be reduced entirely to simple cases by calling the recursive function.
 - *If this is a simple case*
 solve it
 else
 redefine the problem using recursion

Splitting a Problem into Smaller Problems



- Assume that the problem of size 1 can be solved easily (i.e., the simple case).
- We can recursively split the problem into a problem of size 1 and another problem of size $n-1$.

Recursion vs. iteration

- Iteration can be used in place of recursion
 - An iterative algorithm uses a *looping construct*
 - A recursive algorithm uses a *branching structure*
- Recursive solutions are often less efficient, in terms of both *time* and *space*, than iterative solutions
- Recursion can simplify the solution of a problem, often resulting in *shorter*, more easily understood source code

Recursion *vs.* Iteration (Contd...)

- Some simple recursive problems can be “unwound” into loops
 - But code becomes less compact, harder to follow!
- Hard problems cannot easily be expressed in non-recursive code
 - Tower of Hanoi
 - Robots or avatars that “learn”
 - Advanced games

Try it Yourself



- Generate a fibonacci series using recursion

Recursive definition for $\{f_n\}$:

INITIALIZATION: $f_0 = 0, f_1 = 1$

RECURSION : $f_n = f_{n-1} + f_{n-2}$ for $n > 1$

- Finding the GCD
 - Euclid's algorithm makes use of the fact that

$$\gcd(x, y) = \gcd(y, x \bmod y)$$

$$\gcd(x, y) = \begin{cases} x & \text{if } y=0 \\ \gcd(y, x \bmod y) & \text{otherwise} \end{cases}$$

Summary

- Discussed so far the
 - What is a recursion (function) ?
 - What is the need for the recursive function?
 - Writing recursive functions using C
 - How a hard/difficult problem can be solved by recursion
 - Comparison of Recursion with Iterative method